



## Simple Random Walks

### Part V

#### Khinchin's Law of the Iterated Logarithm:

#### Quantitative versions

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In this Fifth Part, we first give a quantitative version of Khinchin's law of the iterated logarithm (1924); we then explain the connections and differences with our present work. The comparison between Khinchin's methods and ours lead to the following conclusions:

- Khinchin's methods may handle the case of a single curve, whereas ours may handle only the case of two barriers;
- Quantitative estimates, about the probability to reach a certain curve at a certain time, obtained by means of Khinchin's methods are of probabilistic type, and are much weaker than similar results derived by means of operator theory;
- To say that Khinchin's curve  $\varphi(x) = \sqrt{2x \log(\log(x))}$  is a "security curve" is incorrect. Take any curve, such as  $b(x) = \sqrt{x \log(x)}$ , which is above Khinchin's curve, and take two instants  $N_1 < N_2$ . Then the probability to hit  $b(x)$  between these two instants, and thus to go above  $k(x)$ , is always strictly positive. The statement about "security curve" is only asymptotic.

# I. Khinchin's Law of the Iterated Logarithm: quantitative version

Let us come back to the original setting: the  $X_n$ 's are independent variables with same law,  $P(X_n = \pm 1) = \frac{1}{2}$ . We set  $S_N = \sum_{n=1}^N X_n$  (in other words, we do not consider only even values of the time and, originally, both fortunes are equal).

We introduce Khinchin's curve, or barrier, defined by the equation:

$$\varphi(x) = \sqrt{2x \text{Log}(\text{Log}(x))},$$

which is a real function, defined for all real  $x > e$ .

The classical statement of the law of the iterated logarithm is:

$$\text{almost surely, } \limsup_{n \rightarrow +\infty} \frac{S_n}{\varphi(n)} = 1$$

(see [https://en.wikipedia.org/wiki/Law\\_of\\_the\\_iterated\\_logarithm](https://en.wikipedia.org/wiki/Law_of_the_iterated_logarithm))

The explanation given by Wikipedia, rather obscure, is as follows:

"Thus, although the quantity  $\left| \frac{S_n}{\varphi(n)} \right|$  is less than any predefined  $\varepsilon > 0$  with probability approaching one, the quantity will nevertheless be greater than  $\varepsilon$  infinitely often; in fact, the quantity will be visiting the neighborhoods of any point in the interval  $(-1,1)$  almost surely."

The difficulty is in the understanding of the words "almost surely", both in theory and in practice.

There is a natural probability on the infinite product  $\prod_{n=1}^{\infty} \{-1,1\}^n$ , which is simply the product of the elementary probabilities on each layer. With respect to this "global" probability, the words "almost surely" are well-defined. But, for this probability, every elementary path has probability 0, and so does a finite number of paths. Moreover, it is quite hard to obtain quantitative results for this probability, which is well-suited only for probabilistic arguments.

On the contrary, if we stop at time  $N$ , we have a precise and intuitive definition of the probability of an event: up to time  $N$ , we have  $2^N$  paths, and the probability of any event is : number of paths satisfying the event, divided by  $2^N$ . For instance, the probability of the event  $\{\exists n \leq N ; S_n > \varphi(n)\}$  is perfectly clear: we count the number of paths for which, at some point, the random walk is above the curve, and we divide this number by  $2^N$ .

Since the paths divide into two at each step, an estimate obtained at a given step will remain valid at later stages. For instance, the statement  $X_1 = 1$  has the same probability (1/2), no matter whether we consider it at stage 1 or at any later stage.

Let us observe that the statement from Wikipedia may be quite misleading. Indeed, if one reads:

"The quantity  $\left| \frac{S_n}{\varphi(n)} \right|$  will nevertheless be greater than  $\varepsilon$  infinitely often", this is true for most curves, and not only for Khinchin's curve. Indeed, we remember (Part I) that the random walk  $S_n$  comes back infinitely many times to the  $x$  axis, so for instance the value  $S_n = 0$  may be expected at time  $n = 5000$ . But then, consider the situation where  $S_n$  increases linearly from this point; it will eventually cross the curve  $\varphi(n)$ ; in fact, this is true for any curve such that  $\frac{\varphi(n)}{n} \rightarrow 0$  when  $n \rightarrow +\infty$ . We construct this way an infinite number of situations in which  $S_n$  exceeds  $4\varphi(n)$  or any multiple, as one wishes.

Therefore, we think that, in such statements, a precise definition of the probabilities must be given. This is what we do now. We give Khinchin's results in quantitative settings, which are new, as far as we know.

**Theorem 1. Quantitative statement of LIL, 1-** *Let  $\varepsilon > 0$  and  $m \geq 2$ . We set:*

$$\eta(m, \varepsilon) = \frac{2}{(\text{Log}(1 + \varepsilon))^{1+\varepsilon}} \frac{1}{\varepsilon(m-1)^\varepsilon}$$

*and:*

$$B(m, \varepsilon) = \left\{ \exists n \geq (1 + \varepsilon)^m, S_n > (1 + \varepsilon)\varphi(n) \right\}$$

*Then, for all  $m$ :*

$$P(B(m, \varepsilon)) \leq \eta(m, \varepsilon).$$

The set  $B(m, \varepsilon)$  is made of the paths which are above the strip  $(1 + \varepsilon)\varphi$  at least once after the time  $(1 + \varepsilon)^m$ . The Theorem says that the probability of such an event tends to 0 when  $m$  increases. In other words, if we fix a width for the strip, that is  $\varepsilon > 0$  fixed, it becomes less and less probable to pass above the curve  $(1 + \varepsilon)\varphi$  when  $n$  increases.

Let us take for example  $\varepsilon = 1$ . Theorem 1 gives the estimate:

$$P\left\{ \exists n \geq 2^m, S_n > 2\varphi(n) \right\} \leq \frac{2}{(\text{Log}(2))^2} \frac{1}{m-1}. \quad (1)$$

For instance, if we want the right-hand side to be  $\leq 0.05$ , we find  $m = 85$ . Therefore, the probability to have  $n \geq 2^{85}$  for which  $S_n > 2\varphi(n)$  is  $< 0.05$ .

For a better understanding, this statement may be converted into a proportion of paths, as follows:

*Fix any  $N > 2^{85}$ . The proportion of paths, which satisfy  $S_n > 2\varphi(n)$  at least in one place, between  $2^{85}$  and  $N$ , is smaller than 5%.*

Still, it is quite possible that a significant number of paths reach a curve above Khinchin's. Let  $b(x)$  be such a curve; an estimate such as:

$$P\{\exists n \geq 2^m, S_n > b(n)\} \leq \frac{c}{2^m} \quad (2)$$

is compatible with (1).

### **Proof of Theorem 1**

Our proof is a quantitative version of the original "Law of the Iterated Logarithm", by A Khinchin. We adapt the presentation given by [Velenik].

In what follows,  $\varepsilon > 0$  (width of the strip) is fixed, so we omit it from most notation. For easy reference, the reader may take  $\varepsilon = 1$ .

We need several steps. We recall from Part I, Lemma 2, that, for all  $n \geq 1$ :

$$P(S_n \geq 0) > \frac{1}{2} \quad (1)$$

We also recall from Part I, Corollary 5, that for any real  $x$  and any  $n \geq 1$ , we have:

$$P(\exists k \leq n, S_k > x) \leq 2P(S_n > x) \quad (2)$$

and, from Part I, Lemma IV.1, that for any  $n$  and any  $x$ ,  $0 \leq x \leq n$ , we have:

$$P(S_n \geq x) \leq e^{\frac{-x^2}{2n}} \quad (3)$$

We define  $\gamma = 1 + \varepsilon$ , and, for all  $k \geq 1$ , integer, we set  $n_k = \gamma^k$ . The next Lemma gives an estimate on the number of paths which are above the strip at least once, in the interval of time  $[n_k, n_{k+1}]$ :

**Lemma 2.** – Let  $C_k = \{\exists n, n_k \leq n < n_{k+1}, S_n > \gamma\varphi(n_k)\}$ . Then:

$$P(C_k) \leq 2(k \text{Log}(\gamma))^{-\gamma}$$

**Proof of Lemma 2**

Using (2), we have:

$$P(C_k) \leq 2P\{S_{n_{k+1}} > \gamma\varphi(n_k)\}$$

and, using (3):

$$P\{S_{n_{k+1}} > \gamma\varphi(n_k)\} \leq \exp\left(-\gamma^2 \frac{n_k}{n_{k+1}} \text{Log}(\text{Log}(n_k))\right)$$

Since  $n_k = \gamma^k$ ,  $\frac{n_{k+1}}{n_k} = \gamma$ , and we get:

$$P(C_k) \leq 2 \exp(-\gamma \text{Log}(\text{Log}(n_k))) = 2(\text{Log}(n_k))^{-\gamma} = 2(k \text{Log}(\gamma))^{-\gamma}$$

This proves Lemma 2.

We set  $D_k = \{\exists n, n_k \leq n < n_{k+1}, S_n > \gamma\varphi(n)\}$ . We have:

**Lemma 3.** – For any  $k$ ,  $D_k \subset C_k$ .

**Proof of Lemma 3**

Indeed, if there exists an  $n$  such that the inequality  $S_n > \gamma\varphi(n)$  holds, we have a fortiori  $S_n > \gamma\varphi(n_k)$ , since the function  $\varphi$  is increasing. This proves Lemma 3. We deduce:

$$P(D_k) \leq 2(k \text{Log}(\gamma))^{-\gamma} \tag{4}$$

Let  $B_m = \bigcup_{k \geq m} D_k$ ; then the sets  $B_m$  are decreasing when  $m$  increases. The set  $B_m$  is by definition:

$$B_m = \bigcup_{k \geq m} \{\exists n, n_k \leq n < n_{k+1}, S_n > \gamma\varphi(n)\} = \{\exists n \geq \gamma^m, S_n > \gamma\varphi(n)\}.$$

Therefore,  $B_m$  is the set of all paths which are above the strip  $\gamma\varphi(n)$  at least once after time  $\gamma^m$ . We now estimate its probability.

**Lemma 4.** – *For all  $m$ , we have:*

$$P(B_m) \leq \frac{2}{(\text{Log } \gamma)^\gamma} \frac{1}{\varepsilon(m-1)^\varepsilon}$$

**Proof of Lemma 4**

By definition of the sets:

$$P(B_m) \leq \sum_{k=m}^{+\infty} P(D_k)$$

and Lemma 4 follows from (4) and the inequality:

$$\sum_{k=m}^{+\infty} \frac{1}{k^{1+\varepsilon}} \leq \int_{m-1}^{+\infty} \frac{dx}{x^{1+\varepsilon}} = \frac{1}{\varepsilon(m-1)^\varepsilon}.$$

The proof of Theorem 1 is complete.

We come back here on what we said about the definition of probabilities. In all the statements above, up to Lemma 5, all probabilities refer to a bounded interval for  $n$  (for instance  $n \leq n_{k+1}$ ). This is not the case for Lemma 6 ( $n \geq \gamma^m$ ), but we immediately have an upper bound from bounded intervals, deduced from (3).

We now turn to the opposite theorem: many paths enter the given strip. More precisely, let  $\alpha > 0$  (small). We will show that there is  $N_0(\varepsilon, \alpha)$  such that if  $N > N_0(\varepsilon, \alpha)$ , then:

$$P\left(\forall n = 1, \dots, N; \frac{S_n}{\varphi(n)} < 1 - \varepsilon\right) < \alpha.$$

The statement is as follows:

**Theorem 5.** – *Let  $\varepsilon > 0, \alpha > 0$ . Set  $\gamma = \frac{16}{\varepsilon^2}$ ,  $\delta = \frac{2}{(\text{Log}(2))^2} \approx 4.2$ ,  $k_0 = 1 + \left(2 + \frac{\delta}{\alpha}\right) \frac{\text{Log}(2)}{\text{Log}(\gamma)}$  and*

*$N_0 = \left(k_0^{\frac{\varepsilon}{2}} + \frac{\varepsilon}{2} \text{Log}\left(\frac{2}{\alpha}\right)\right)^{\frac{2}{\varepsilon}}$ . For any  $N \geq N_0$ , we have:*

$$P\left(\exists k, k_0 \leq k \leq N, \text{ such that } \frac{S_k}{\varphi(k)} > 1 - \varepsilon\right) \geq 1 - \alpha.$$

What this statement says, in simple words, is that it is more and more unlikely to stay constantly below  $(1 - \varepsilon)\varphi(x)$ . For a fixed width  $\varepsilon$ , the probability that  $S_n$  enters, at some time, the strip  $[(1 - \varepsilon)\varphi(x), \varphi(x)]$  tends to 1 when  $n \rightarrow +\infty$ .

### Proof of Theorem 5

We set as before  $\varphi(n) = \sqrt{2n \text{Log}(\text{Log}(n))}$ . Let  $\varepsilon > 0$  : it refers to the width of the strip, is fixed and is most of the time omitted from the notation. We introduce  $\gamma = \frac{16}{\varepsilon^2}$  and, for any integer  $k$ ,  $n_k = \gamma^k$ . We write the quotient  $\frac{S_{n_k}}{\varphi(n_k)}$  under the form of the sum of two terms, which will be treated separately:

$$\frac{S_{n_k}}{\varphi(n_k)} = \frac{S_{n_k} - S_{n_{k-1}}}{\varphi(n_k)} + \frac{S_{n_{k-1}}}{\varphi(n_{k-1})} \frac{\varphi(n_{k-1})}{\varphi(n_k)} \quad (1)$$

Set  $Y_k = \frac{S_{n_k} - S_{n_{k-1}}}{\varphi(n_k)}$  and  $Z_k = \frac{S_{n_{k-1}}}{\varphi(n_{k-1})} \frac{\varphi(n_{k-1})}{\varphi(n_k)}$ . The general idea of the proof is to show that  $Y_k$  is large with large probability, whereas  $Z_k$  is small. We start with the study of  $Z_k$ .

**Lemma 6.** – For all  $k \geq 2$ , we have:

$$\frac{\varphi(n_{k-1})}{\varphi(n_k)} \leq \frac{1}{\sqrt{\gamma}} = \frac{\varepsilon}{4}$$

### Proof of Lemma 6

This is clear, since  $\frac{\varphi(n_{k-1})}{\varphi(n_k)} = \frac{\sqrt{\gamma^{k-1} \text{Log}((k-1) \text{Log}(\gamma))}}{\sqrt{\gamma^k \text{Log}(k \text{Log}(\gamma))}}$  and by the choice of  $\gamma$ .

**Lemma 7.** – Let  $k_0 = 1 + \left(\frac{4}{(\text{Log}(2))^2} \frac{1}{\alpha} + 1\right) \frac{\text{Log}(2)}{\text{Log}(\gamma)}$ . Then:

$$P\left(\forall k \geq k_0, \frac{|S_{n_{k-1}}|}{\varphi(n_{k-1})} < 2\right) \geq 1 - \frac{\alpha}{2}.$$

### Proof of Lemma 7

Let us choose  $\varepsilon = 1$  in Theorem 1. We have, for any  $m$  :

$$P\left(\left\{\exists n \geq 2^m, S_n > 2\varphi(n)\right\}\right) \leq \frac{2}{(\text{Log}(2))^2} \frac{1}{m-1}$$

We choose  $m$  so that  $\frac{2}{(\text{Log}(2))^2} \frac{1}{m-1} < \frac{\alpha}{2}$ , that is  $m > \frac{4}{\text{Log}^2 2} \frac{1}{\alpha} + 1$ . So we get:

$$P\left(\exists n > 2^m, |S_n| > 2\varphi(n)\right) < \frac{\alpha}{2}$$

that is:

$$P\left(\forall n > 2^m, \frac{|S_n|}{\varphi(n)} \leq 2\right) \geq 1 - \frac{\alpha}{2}$$

Here,  $n = \gamma^{k-1}$  and the condition  $n \geq 2^m$  is satisfied as soon as:

$$k > 1 + m \frac{\text{Log}(2)}{\text{Log}(\gamma)} > 1 + \left( \frac{4}{(\text{Log}(2))^2} \frac{1}{\alpha} + 1 \right) \frac{\text{Log}(2)}{\text{Log}(\gamma)}$$

This proves Lemma 7.

**Lemma 8.** - Let  $k_0$  be as before. If  $k \geq k_0$ , we have:

$$P\left(\forall k \geq k_0, |Z_k| < \frac{\varepsilon}{2}\right) \geq 1 - \frac{\alpha}{2}.$$

### Proof of Lemma 8

Indeed, for any  $k$ , we have both  $\frac{|S_{n_{k-1}}|}{\varphi(n_{k-1})} < 2$  and  $\frac{\varphi(n_{k-1})}{\varphi(n_k)} < \frac{\varepsilon}{4}$  on a set of probability  $\geq 1 - \frac{\alpha}{2}$ .

This proves Lemma 8.

We now turn to the term  $Y_k = \frac{S_{n_k} - S_{n_{k-1}}}{\varphi(n_k)}$ . We recall from Part I, Proposition 9 that, if  $k < \sqrt{n}$ , we

have, with  $c = \frac{1}{4\sqrt{2\pi}}$ :



$$P(S_n \geq k) \geq c \exp\left(-\frac{k^2}{2n}\right) \quad (2)$$

We set:

$$D_k = \left\{ \frac{S_{n_k} - S_{n_{k-1}}}{\varphi(n_k)} > 1 - \frac{\varepsilon}{2} \right\}$$

**Proposition 9.** – *For all  $k \geq 1$ , we have:*

$$P(D_k) \geq \frac{c}{(k \text{Log}(\gamma))^{1-\varepsilon/2}},$$

where  $c = \frac{1}{4\sqrt{2\pi}}$  as before.

### Proof of Proposition 9

Since the definition of  $D_k$  relies upon consecutive differences  $S_{n_k} - S_{n_{k-1}}$ , the events  $D_k$  are independent. We have:

$$P(D_k) = P\left\{ S_{n_k} - S_{n_{k-1}} > \left(1 - \frac{\varepsilon}{2}\right) \varphi(n_k) \right\} = P\left\{ S_{n_k - n_{k-1}} > \left(1 - \frac{\varepsilon}{2}\right) \varphi(n_k) \right\}$$

since  $S_{n_k} - S_{n_{k-1}}$  and  $S_{n_k - n_{k-1}}$  have the same law.

In the estimate (2) above, we replace  $n$  by  $n_k - n_{k-1}$  and  $k$  by  $\left(1 - \frac{\varepsilon}{2}\right) \varphi(n_k)$ ; we obtain:

$$P\left( S_{n_k - n_{k-1}} \geq \left(1 - \frac{\varepsilon}{2}\right) \varphi(n_k) \right) \geq c \exp\left( -\frac{1}{2} \left(1 - \frac{\varepsilon}{2}\right)^2 \frac{(\varphi(n_k))^2}{n_k - n_{k-1}} \right) = c \exp\left( -\left(1 - \frac{\varepsilon}{2}\right)^2 \frac{n_k \text{Log}(\text{Log}(n_k))}{n_k - n_{k-1}} \right)$$

But:

$$\frac{n_k - n_{k-1}}{n_k} = \frac{\gamma^k - \gamma^{k-1}}{\gamma^k} = \frac{\gamma - 1}{\gamma} > 1 - \frac{\varepsilon}{2}, \text{ from the choice of } \gamma. \text{ Therefore:}$$

$$P\left\{ S_{n_k - n_{k-1}} > \left(1 - \frac{\varepsilon}{2}\right) \varphi(n_k) \right\} \geq c \exp\left( -\left(1 - \frac{\varepsilon}{2}\right) \text{Log}(\text{Log}(n_k)) \right) = c (\text{Log}(n_k))^{-\left(1 - \frac{\varepsilon}{2}\right)} = \frac{c}{(k \text{Log}(\gamma))^{1-\varepsilon/2}}$$

This proves Proposition 9.

We need a quantitative version of the second Borel-Cantelli Lemma:

**Lemma 10.** – *Let  $D_k$  be a sequence of independent events; let  $D_k^c$  be their complements. Then :*

$$P\left(\bigcap_{k=1}^N D_k^c\right) \leq \exp\left(-\sum_{k=1}^N P(D_k)\right)$$

**Proof of Lemma 10**

We have, for all  $N$  :

$$P\left(\bigcap_{k=1}^N D_k^c\right) = \prod_{k=1}^N (1 - P(D_k)) \leq \exp\left(-\sum_{k=1}^N P(D_k)\right)$$

using the inequality  $1 - x \leq e^{-x}$ ,  $0 \leq x \leq 1$ . This proves Lemma 10.

Set, for all  $k$ ,  $u_k = \frac{c}{(\text{Log}(\gamma))^{1-\varepsilon/2}} \frac{1}{k^{1-\varepsilon/2}}$ . We deduce from Proposition 9 and Lemma 10:

$$P\left(\bigcap_{k=1}^N D_k^c\right) \leq \exp\left(-\sum_{k=1}^N u_k\right).$$

Since the series of general term  $u_k$  is divergent, the sum  $\sum_{k=1}^N u_k$  can be made arbitrarily large, choosing  $N$  large enough. Then  $\exp\left(-\sum_{k=1}^N u_k\right)$  is close to 0. So the intersection  $\bigcap_{k=1}^N D_k^c$  has a very small probability. But this intersection is the set of all paths for which  $\left\{ \frac{S_{n_k} - S_{n_{k-1}}}{\varphi(n_k)} \leq 1 - \frac{\varepsilon}{2} \right\}$  for any  $k = 1, \dots, N$ . More precisely, we have:

**Lemma 11.** – *Let  $k_0$  as in Lemma 7. We have, if  $N > N_0 = \left( k_0^{\frac{\varepsilon}{2}} + \frac{\varepsilon}{2} \text{Log}\left(\frac{2}{\alpha}\right) \right)^{\frac{2}{\varepsilon}}$  :*

$$P\left(\forall k, k_0 \leq k \leq N, \frac{S_{n_k} - S_{n_{k-1}}}{\varphi(n_k)} \leq 1 - \frac{\varepsilon}{2}\right) < \frac{\alpha}{2}.$$

**Proof of Lemma 11**

We choose  $N_0$  large enough so that, if  $N > N_0$ ,  $\exp\left(-\sum_{k=k_0}^N u_k\right) < \frac{\alpha}{2}$ , that is  $\sum_{k_0}^N u_k > \text{Log}\left(\frac{2}{\alpha}\right)$ .

But, for  $k > k_0$ , we have  $u_k > \frac{1}{k^{1-\frac{\varepsilon}{2}}}$ , and therefore:

$$\sum_{k_0}^N u_k > \sum_{k_0}^N k^{-1+\frac{\varepsilon}{2}} \geq \frac{2}{\varepsilon} \left( N^{\frac{\varepsilon}{2}} - k_0^{\frac{\varepsilon}{2}} \right).$$

So we choose  $N$  large enough, in order to have:

$$\frac{2}{\varepsilon} \left( N^{\frac{\varepsilon}{2}} - k_0^{\frac{\varepsilon}{2}} \right) > \text{Log}\left(\frac{2}{\alpha}\right)$$

that is:

$$N \geq \left( k_0^{\frac{\varepsilon}{2}} + \frac{\varepsilon}{2} \text{Log}\left(\frac{2}{\alpha}\right) \right)^{\frac{2}{\varepsilon}}$$

This proves Lemma 11.

Let us finish the proof of Theorem 5. We have:

$$\frac{S_{n_k}}{\varphi(n_k)} = \frac{S_{n_k} - S_{n_{k-1}}}{\varphi(n_k)} + \frac{S_{n_{k-1}}}{\varphi(n_{k-1})} \frac{\varphi(n_{k-1})}{\varphi(n_k)} \quad (3)$$

Let  $E_1$  be the set:

$$E_1 = \left\{ Y_k < 1 - \frac{\varepsilon}{2}, \forall k = k_0, \dots, N \right\}$$

Then, Lemma 11 says that  $P(E_1) < \frac{\alpha}{2}$ .

Let  $E_2$  be the set:

$$E_2 = \left\{ \forall k \geq k_0, |Z_k| < \frac{\varepsilon}{2} \right\}$$

Then  $P(E_2) > 1 - \frac{\alpha}{2}$ . We have:

$$P\left(\left(E_1^c \cap E_2\right)^c\right) = P\left(E_1 \cup E_2^c\right) \leq P\left(E_1\right) + P\left(E_2^c\right) = P\left(E_1\right) + 1 - P\left(E_2\right) < \frac{\alpha}{2} + \frac{\alpha}{2} = \alpha$$

and therefore:

$$P\left(E_1^c \cap E_2\right) \geq 1 - \alpha$$

But  $E_1^c \cap E_2$  is the set for which there is a  $k$ ,  $k_0 \leq k \leq N$ , with  $Y_k \geq 1 - \frac{\varepsilon}{2}$  and  $|Z_k| < \frac{\varepsilon}{2}$ . We deduce:

$$Y_k + Z_k > Y_k - |Z_k| > 1 - \varepsilon$$

This finishes the proof of Theorem 5.

## II. Comparison with our results

If we consider  $\varphi(x) = \sqrt{2x \operatorname{Log}(\operatorname{Log}(x))}$  then the integral  $\int \frac{1}{\varphi^2}$  diverges at  $x = +\infty$ . The same holds for the barrier  $b(x) = \sqrt{x \operatorname{Log}(x)}$ , which is above the previous one. We know that:

$$\int \frac{1}{b^2(x)} dx = \int \frac{1}{x \operatorname{Ln}(x)} dx = \operatorname{Ln}(\operatorname{Ln}(x))$$

and therefore the integral diverges at  $+\infty$ . So, the total energy left at time  $N$  (probability that the game continues up to time  $N$ ) tends to zero when  $N \rightarrow +\infty$ , for both barriers.

The comparison between Khinchin's methods and ours lead to the following conclusions:

- Khinchin's methods may handle the case of a single curve, whereas ours may handle only the case of two symmetric barriers;
- Quantitative estimates obtained by means of Khinchin's methods are of probabilistic type, and much weaker than the results derived by means of operator theory;
- To say that Khinchin's curve  $\varphi(x) = \sqrt{2x \operatorname{Log}(\operatorname{Log}(x))}$  is a "security curve" is incorrect ; the statement *almost surely*,  $\limsup_{n \rightarrow +\infty} \frac{S_n}{\varphi(n)} = 1$  is only asymptotic.

## References

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